



دانشگاه صنعتی اصفهان
دانشکده مهندسی حمل و نقل

حمل و نقل ریلی

مسئله گروه بندی واگن ها

Railroad Blocking Problem

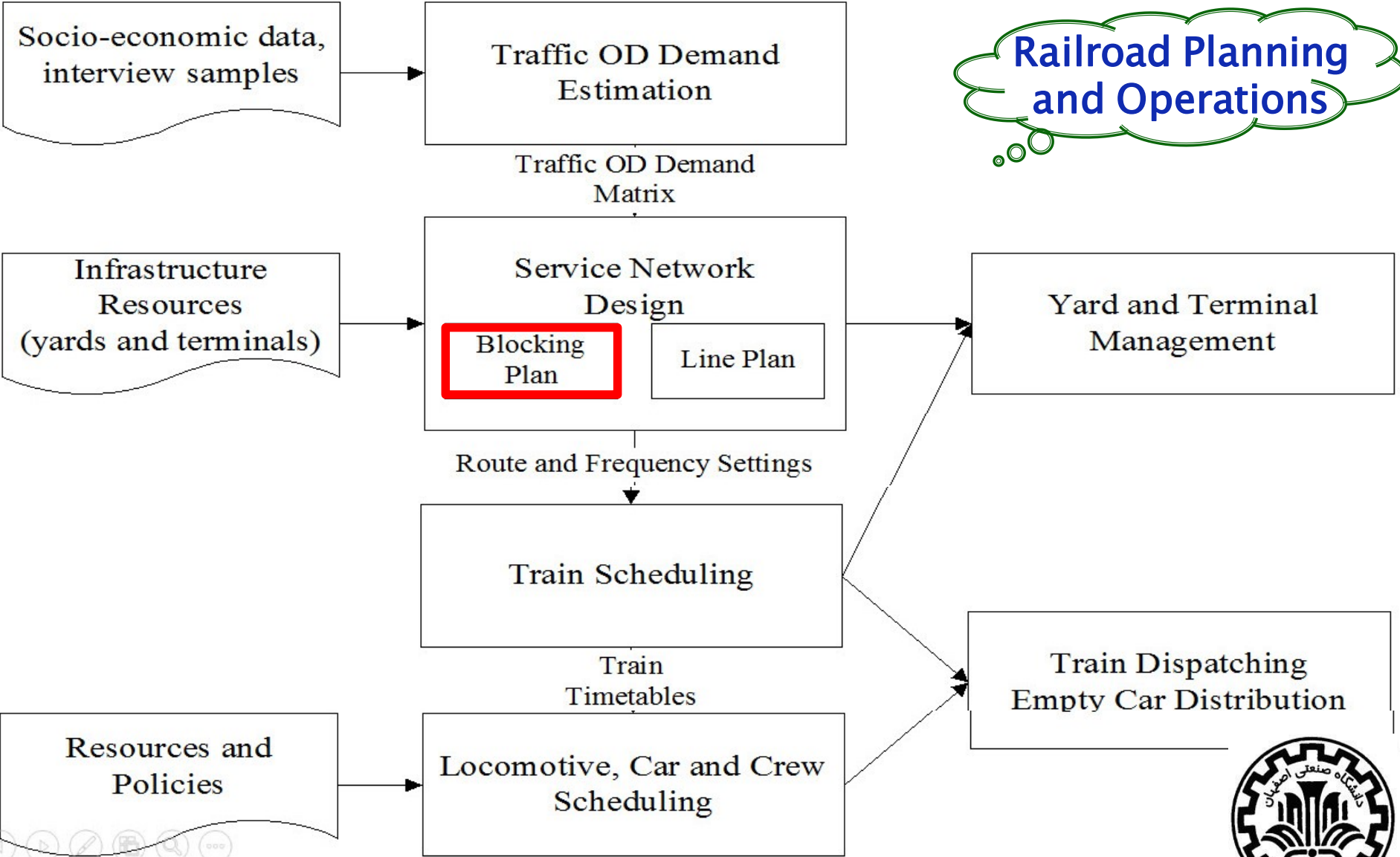
مدرس: محمد تمنایی

بهار ۱۳۹۴

Railroad Blocking Problem:

- ✓ Introduction
- ✓ Significance
- ✓ Terminology
- ✓ Description
- ✓ Example
- ✓ Mathematical model
- ✓ References





Railroad Planning and Operations

OD Demand --> Routes --> Blocks --> Trains

- ▶ Railroad planning and scheduling problems are very large-scale and very difficult discrete optimization problems.
- ▶ In large U.S. Networks, there are teams of 10-20 highly experienced personnel for solving each problem.



Introduction

Railroad Blocking Problem (RBP)



حمل و نقل ریلی

Railroad Planning and Operations



وزارت راه و شهرسازی
راه آهن جمهوری اسلامی ایران

و فناوری اطلاعات
نقشه بر مبنای
GIS

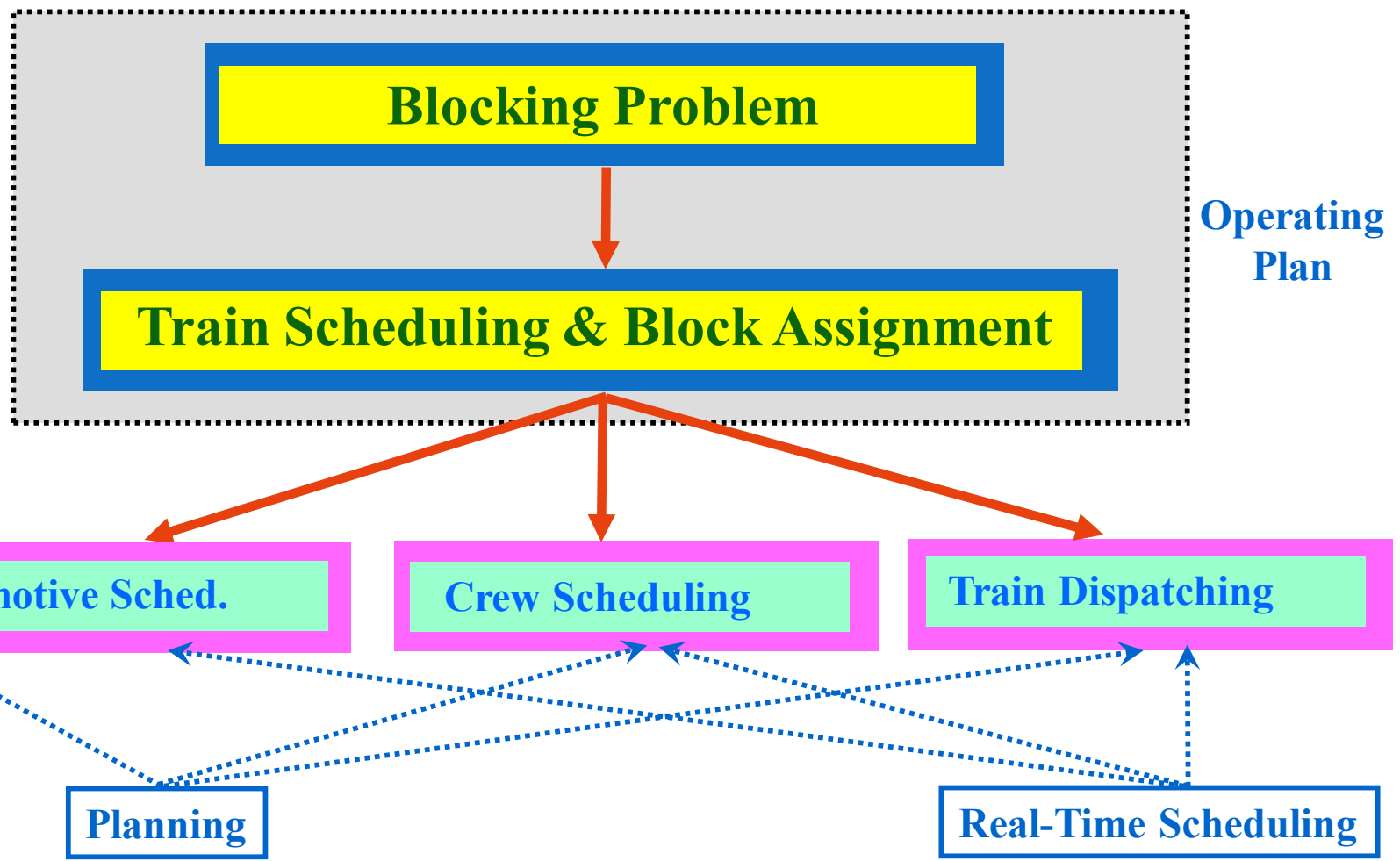
OD Demand --> Routes --> Blocks --> Trains

The questions which must be answered:

- ▶ **Origin Destination (OD) Traffic Demand**
 - Commodities, commercial unit traffic groups, car equipment requirement ?
- ▶ **Route Plan**
 - Single route or multiple routes for each OD pair ?
- ▶ **Blocking Plan**
 - Block origin, block destination, commodities, final destinations ?
- ▶ **Make-up Plan**
 - Block-to-train assignment ?
- ▶ **Train Schedule**
 - Arrival and departure times, station dwell times ?
 - order of departure?

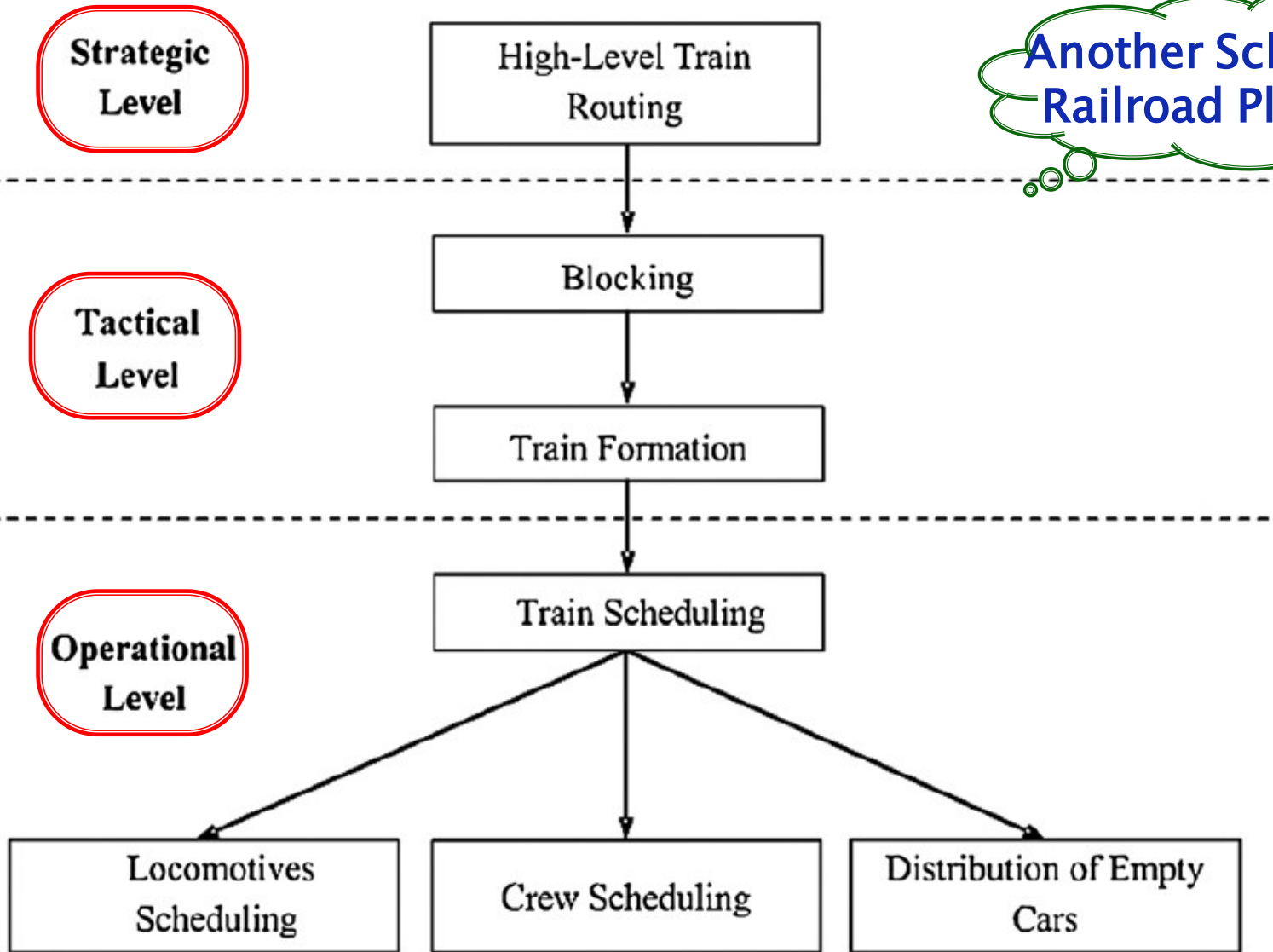


Another Scheme of Railroad Planning



Another Scheme of Railroad Planning

Hierarchical planning in demand-driven railways



What is “Blocking problem” ?

Problem of determining how to aggregate a large number of commodities (shipments) into blocks of commodities as they travel from origins to destinations.

✓ **Blocking problem:**

Assignment of shipments to blocks

And

Routing of blocks through pre-determined paths in the network.

✓ **Train Formation (Make-up) Problem:**

✓ Assignment of blocks to specific trains.

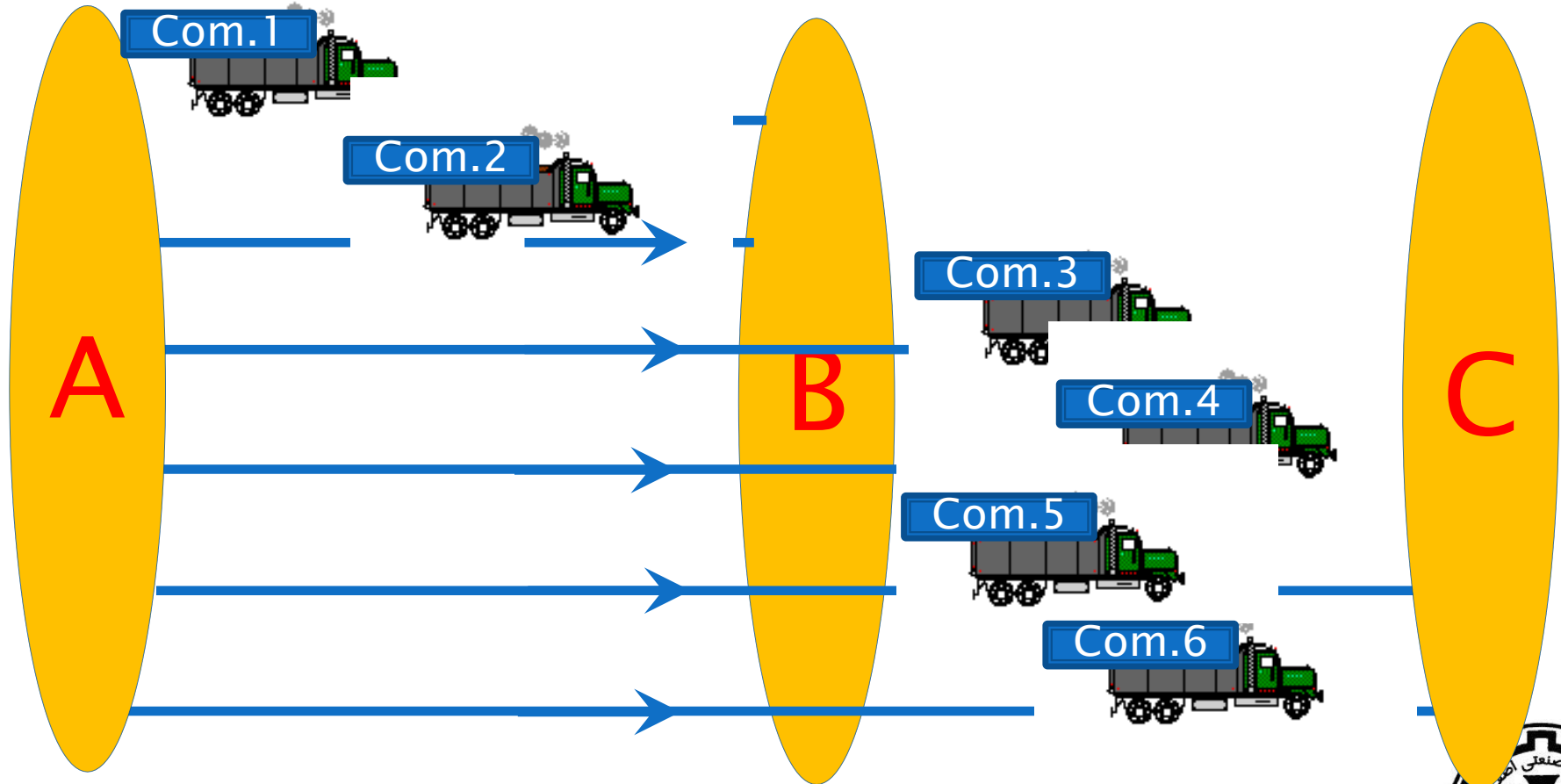
✓ Next step after developing Blocking plan:

To determine in which trains those blocks should travel (Train Formation Problem).



6 Commodities should be transferred:

- | | | |
|-------|-------|-------|
| Com.1 | Com.3 | Com.5 |
| Com.2 | Com.4 | Com.6 |

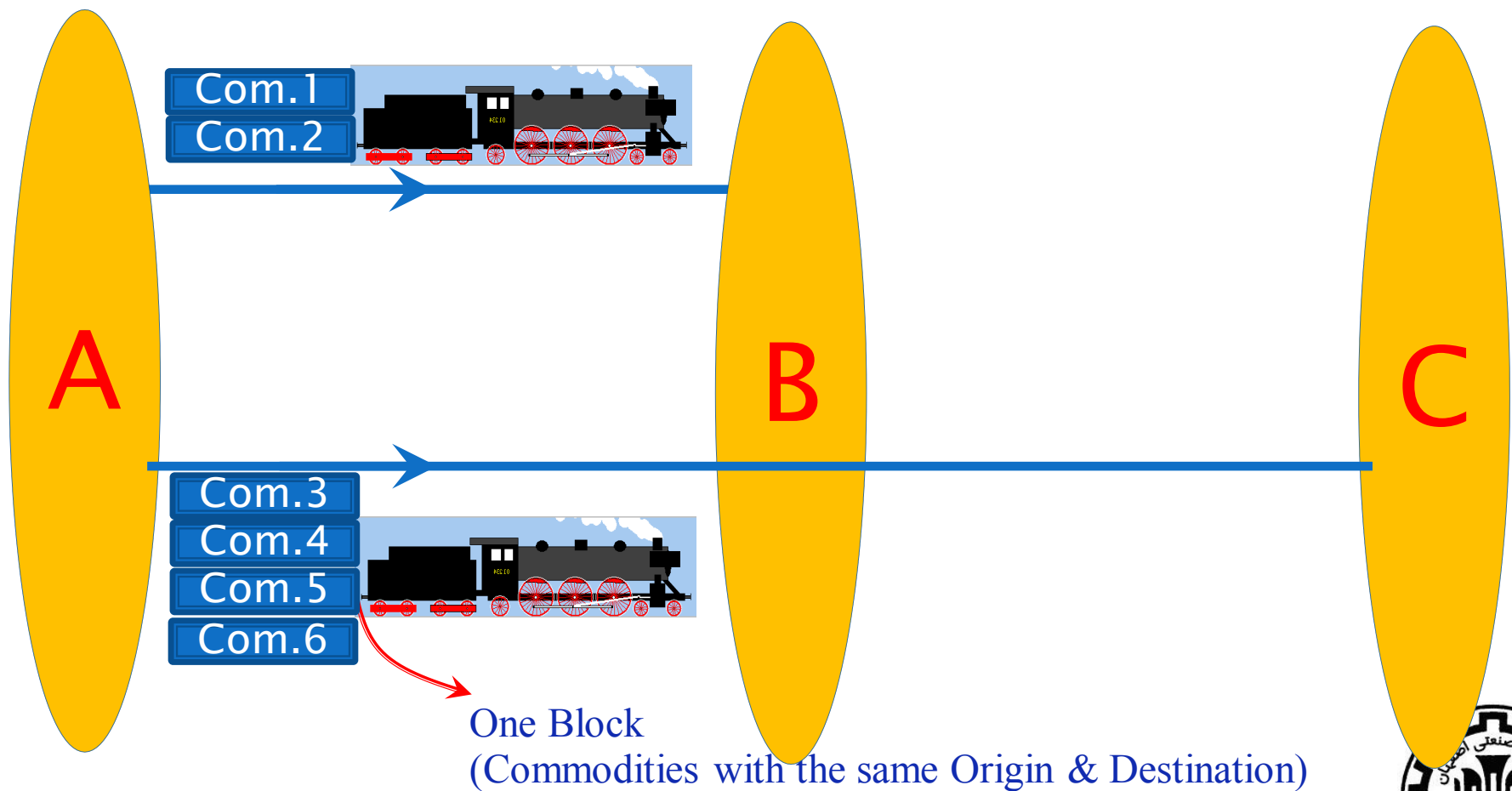


Transportation in Road System:
Each commodity is transferred by one truck.



6 Commodities should be transferred:

- Com.1
- Com.2
- Com.3
- Com.4
- Com.5
- Com.6



Transportation in Railroad System:
Commodities should be aggregated as blocks.



- Com.1
- Com.2
- Com.3
- Com.4
- Com.5
- Com.6

How to aggregate ?

Block 1

Block 2

Block 3

<div data-bbox="850 399 1081 542">Com.1 Com.2</div>	<div data-bbox="1217 314 1449 599">Com.3 Com.4 Com.5 Com.6</div>	----
<div data-bbox="850 742 1081 956">Com.1 Com.2 Com.3</div>	<div data-bbox="1217 742 1449 956">Com.4 Com.5 Com.6</div>	----
<div data-bbox="850 1142 1081 1285">Com.1 Com.2</div>	<div data-bbox="1217 1142 1449 1285">Com.3 Com.4</div>	<div data-bbox="1603 1142 1835 1285">Com.5 Com.6</div>



- ✓ **Goal of Blocking plan:** to minimize the costs.
- ✓ **Two distinct costs:** distance-traveled costs, handling costs.
- ✓ **Distance-traveled costs:** cost arising from various blocks moving through railway network
- ✓ **Handling costs:** related to intermediate reclassification of shipments at classification yards
- ✓ **Hump yards** have lower handling costs.
- ✓ **Flat yards** have higher handling costs.



- ✓ An efficient blocking plan can reduce total operating costs in railroad operations. Classifications at yards are labor and capital-intensive, consisting of 10% of railroad total operating cost on average.
- ✓ The blocking plan has ripple effects on subsequent plans, including train scheduling, crew and power assignment that are developed based on given blocking plans.
- ✓ A good blocking plan has the potential to improve railroad service levels. Through reductions in number of classifications, a good blocking plan can decrease the potential delays occurring in classifications yards, thereby enhancing service quality and, in turn, improving the ability of the railroad to compete with other freight transportation modes, such as trucking and airlines.



Difficulty of the Problem

- ✓ Mathematically, the railroad blocking problem is a multicommodity-flow, network-design, and routing problem.
- ✓ Network-design problems are from NP-Hard problems.

$$\text{Minimize } z(f, y) = \sum_{k \in K} \sum_{(i, j) \in A} c_{ij}^k f_{ij}^k + \sum_{(i, j) \in A} F_{ij} y_{ij}$$

- ✓ Problems with only a few hundred network design variables can be solved to optimality.
- ✓ Railroads want a near-optimal and implementable solution within a few hours of computational time.
- ✓ Real-life blocking problems often contain over a million design variables (F_{ij}) and hundreds of billions of flow variables (f_{ij}^k).



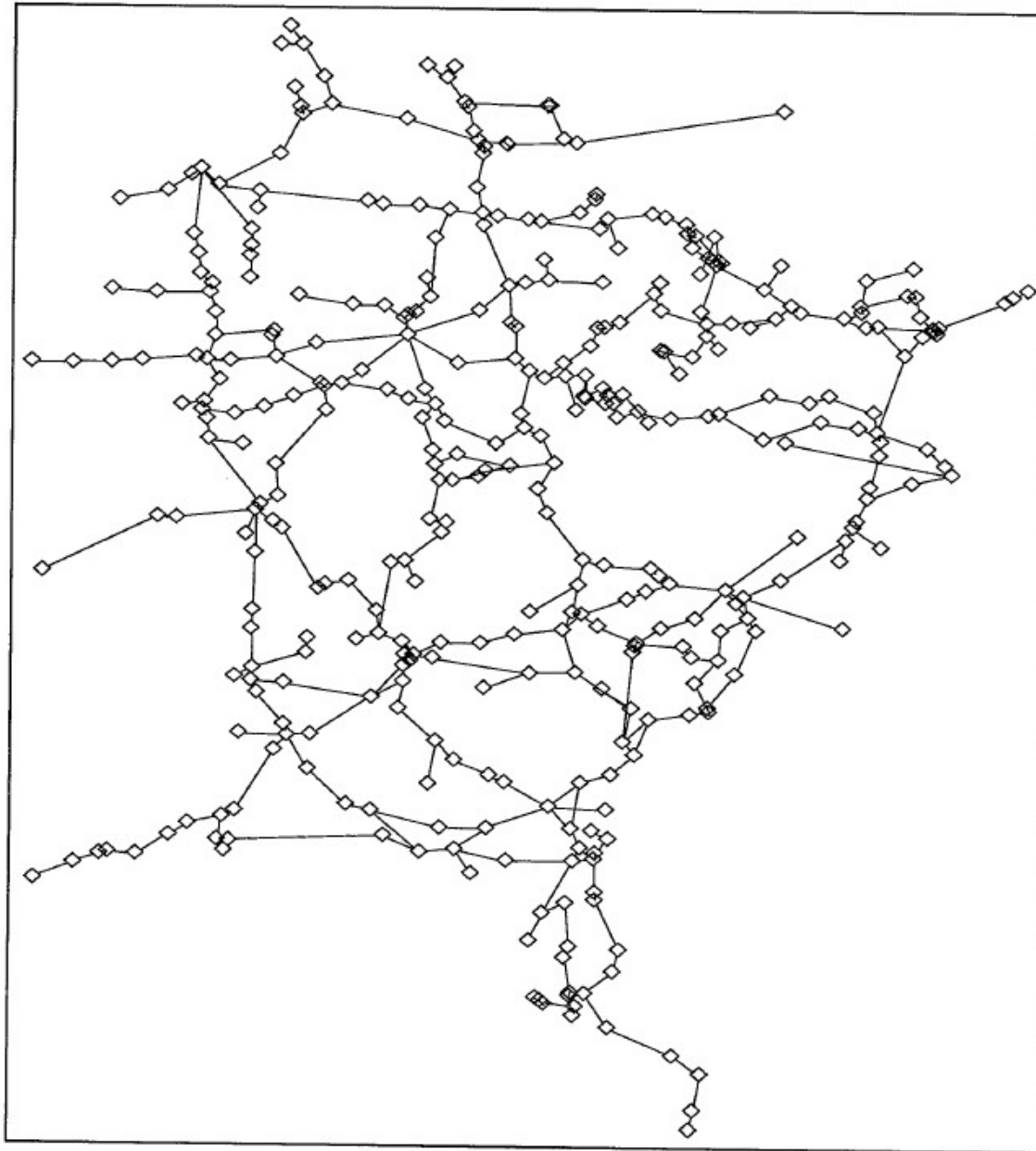


Figure 3.2: Physical rail network of links and nodes



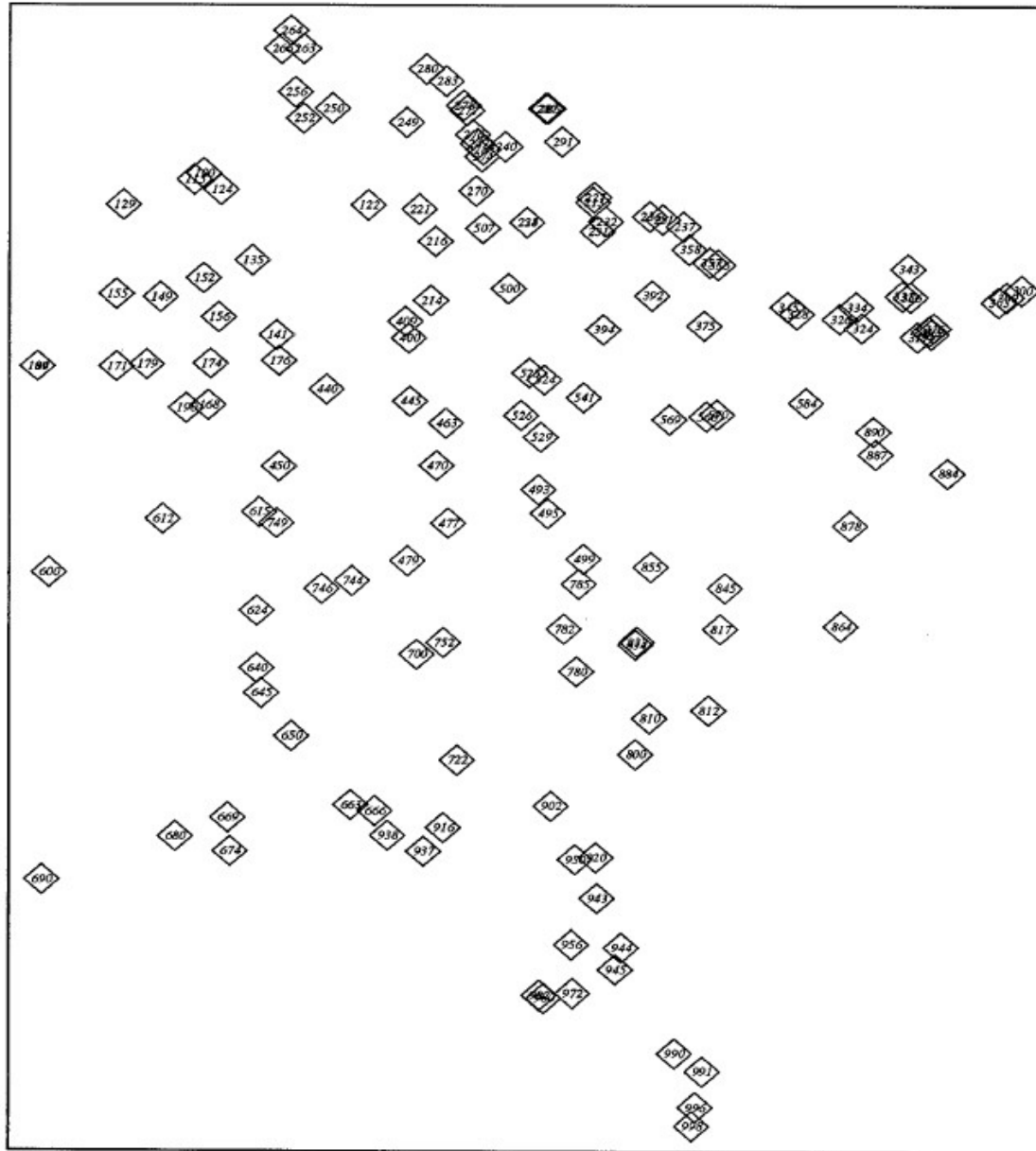


Figure 3.3: Locations of terminals in the rail network



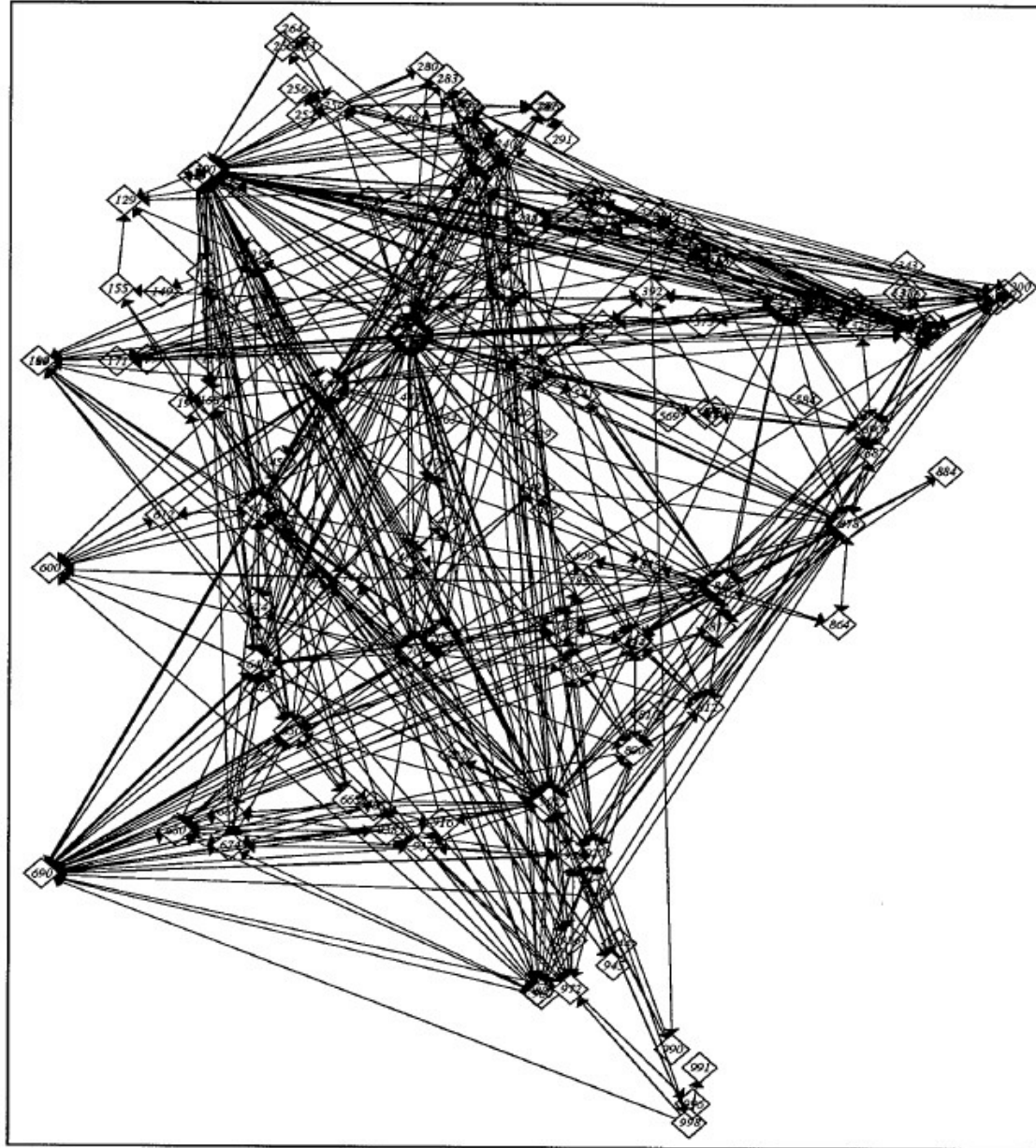


Figure 4.1: Solution for RBP for rail system of Figures 3.2 and 3.3



Block:

A Block is completely specified by its origin and destination.

Cars using the block will be sorted at the block's origin and then not sorted again until reaching the block's destination.

Destination of a block is not necessarily equal to destination(s) of its cars.

Blocking:

Classification, Blocking, or Handling is the process of sorting cars into different blocks.

Terminal:

A node of the rail network at which classification may be performed.



OD pair:

A pair of Origin- Destination, which is the same for group of cars, put in a block.

Priority Class:

Number of intermediate classifications that are permitted.

If no intermediate classifications are permitted, then the commodity must be blocked for the destination terminal directly at the origin terminal.

Commodity (Shipment):

a subgroup of an OD pair demand which has the same priority class.

Routing for a commodity:

Consists of a sequence of terminals, starting with the commodity's origin and ending with the commodity's destination, that a commodity may visit.

A Train:

consists of one or more blocks, in addition to the locomotives.



Terminology

Railroad Blocking Problem (RBP)

Classification yard (Station Grade 1):

A railway yard used to separate, sort and group cars according to their final destination(s), and also allows the inspection of trains.

Classification Tracks

Arrival Tracks

Name	Arr Time
------	----------

BOOHON367,BOONE_VA_Inspecting
EMPHON360,EMPORIA_KS_Waiting For Hump
SOHHON373,SOMERSET_VA_Inspecting
CHEHON071,CHEMULT_OB_Wait for Inspection
ROHHON333,ROBBINSVILLE_Wait for Inspection
CHEHON072,CHEMULT_OB_Wait for Inspection

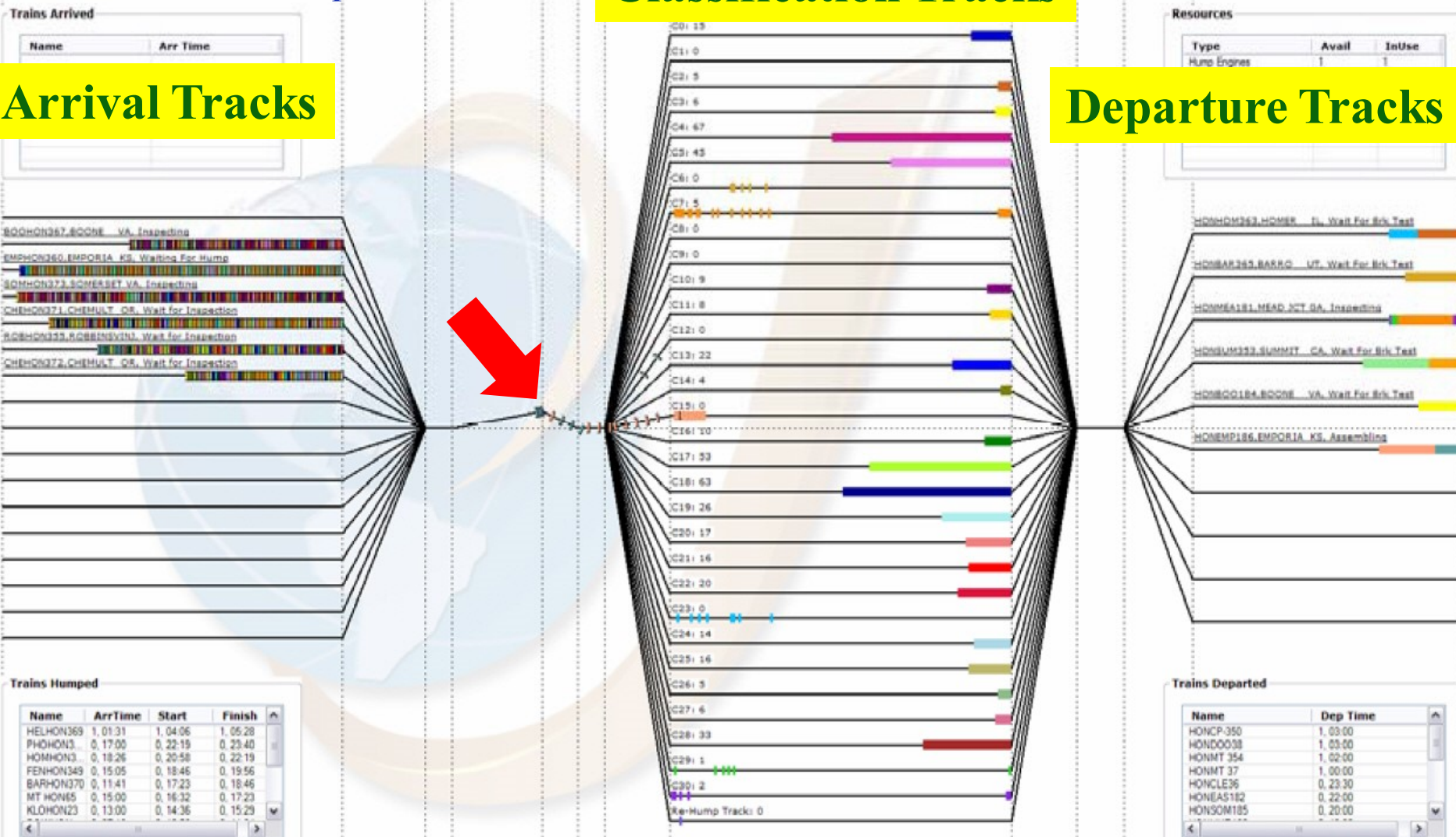
Name	ArrTime	Start	Finish
HELHON369	1, 01:31	1, 04:06	1, 05:28
PHOHON3	0, 17:00	0, 22:19	0, 23:40
HOMHON3	0, 18:26	0, 20:58	0, 22:19
FENHON349	0, 15:05	0, 18:46	0, 19:56
BARHON370	0, 11:41	0, 17:23	0, 18:46
MT HON65	0, 15:00	0, 16:32	0, 17:23
KLOHON23	0, 13:00	0, 14:36	0, 15:29

Departure Tracks

Type	Avail	InUse
Hump Engines	1	1

HONHON363,HOMER_IL_Wait For Bk Test
HONBAR363,BARRON_VT_Wait For Bk Test
HONME481,MEAD_NCT_GA_Inspecting
HONSUM333,SUMMIT_CA_Wait For Bk Test
HONBO0184,BOONE_VA_Wait For Bk Test
HONEMP186,EMPORIA_KS_Assembling

Name	Dep Time
HONCP 350	1, 03:00
HONDO038	1, 03:00
HONMT 354	1, 02:00
HONMT 37	1, 00:00
HONCLE36	0, 23:30
HONEA5182	0, 22:00
HONSON185	0, 20:00

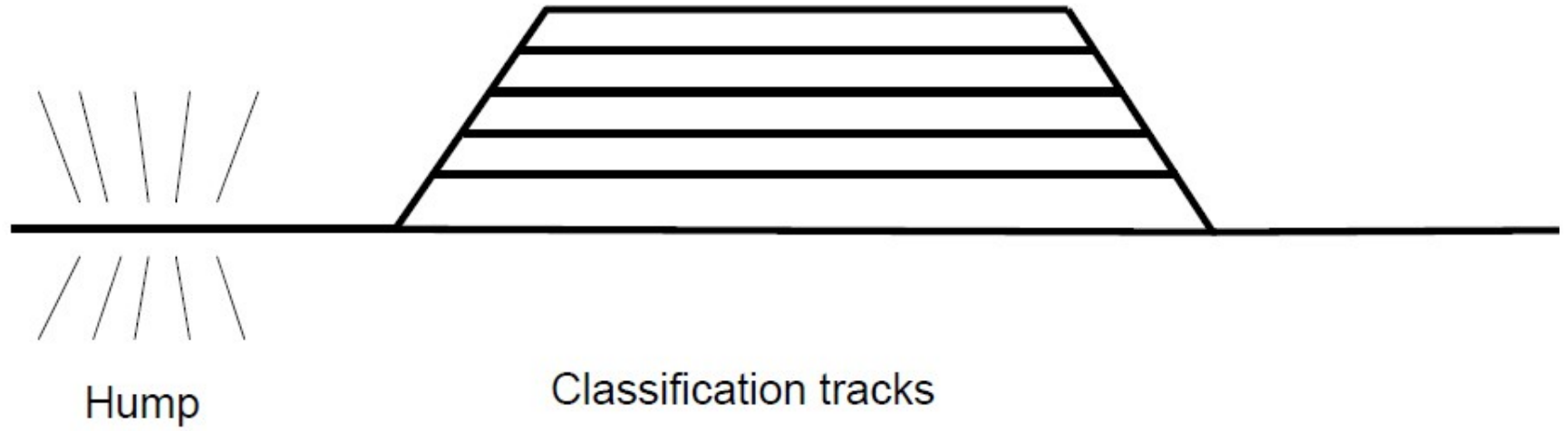


Classification yard (Station Grade 1):

- ✓ **There are 2 types of classification yards:** hump yards and flat yards.
- ✓ **Hump yards** are characterized by a “hump” or hill, upon which the cars are pushed by an engine and moved down by gravity.
- ✓ **At the top of the hill,** cars are automatically decoupled and switched to the proper destination track, onto which, via gravity, they roll down and are now grouped with other cars heading on the same destination track.
- ✓ **Hump yards** are the most cost-effective and with the largest capacity.
- ✓ **In flat yards,** the cars are moved onto the designated tracks by a locomotive, rather than via gravity assist.



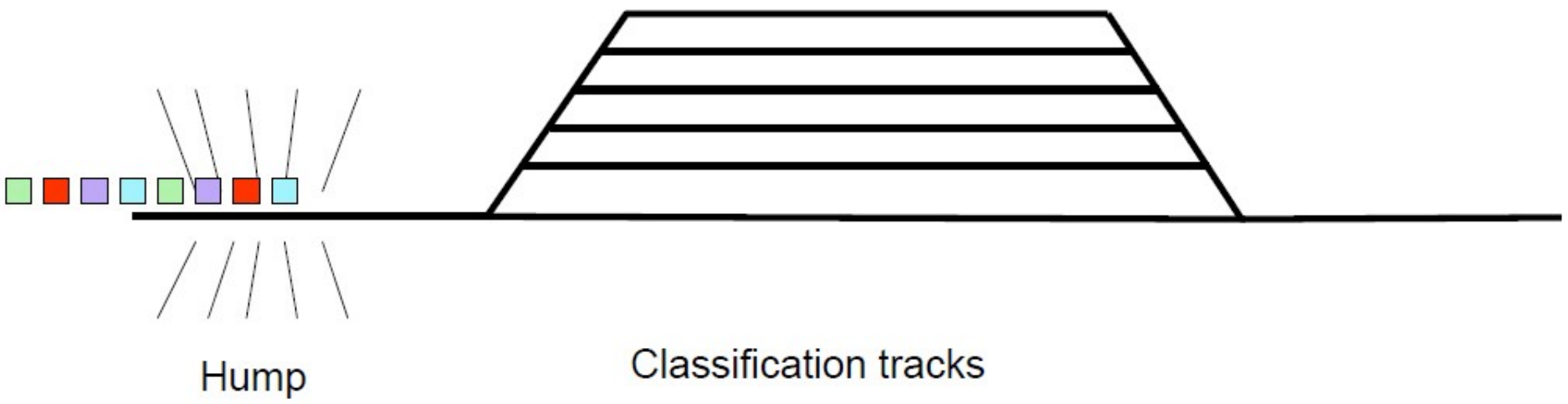
Classification yard



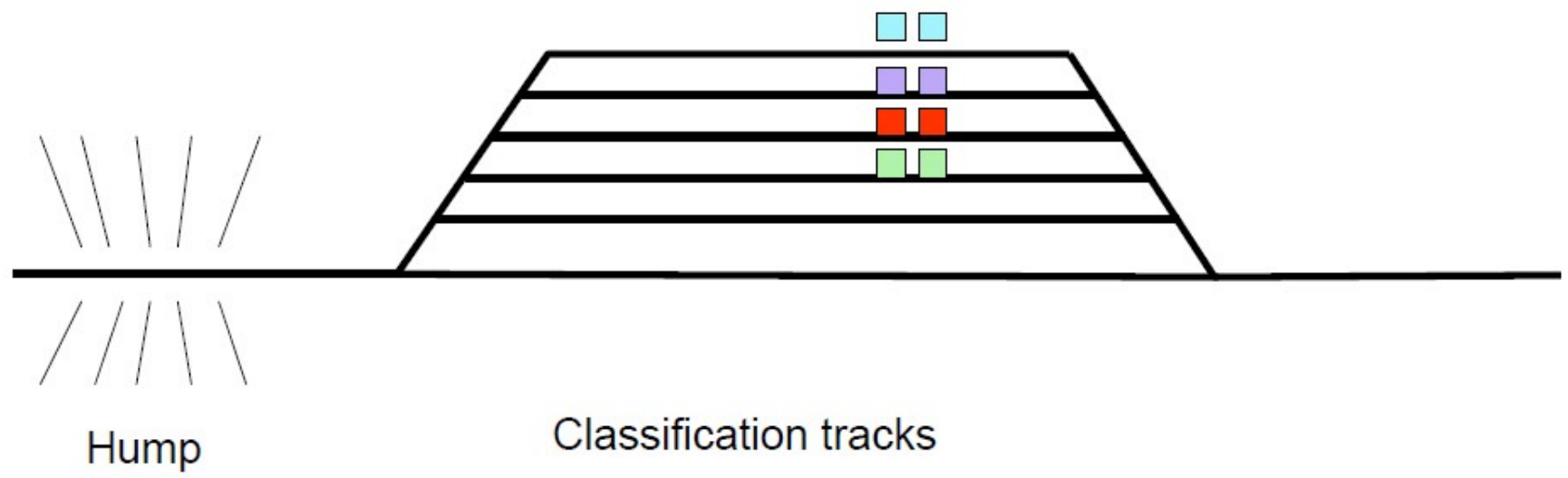
The commodities are reclassified in “classification tracks”



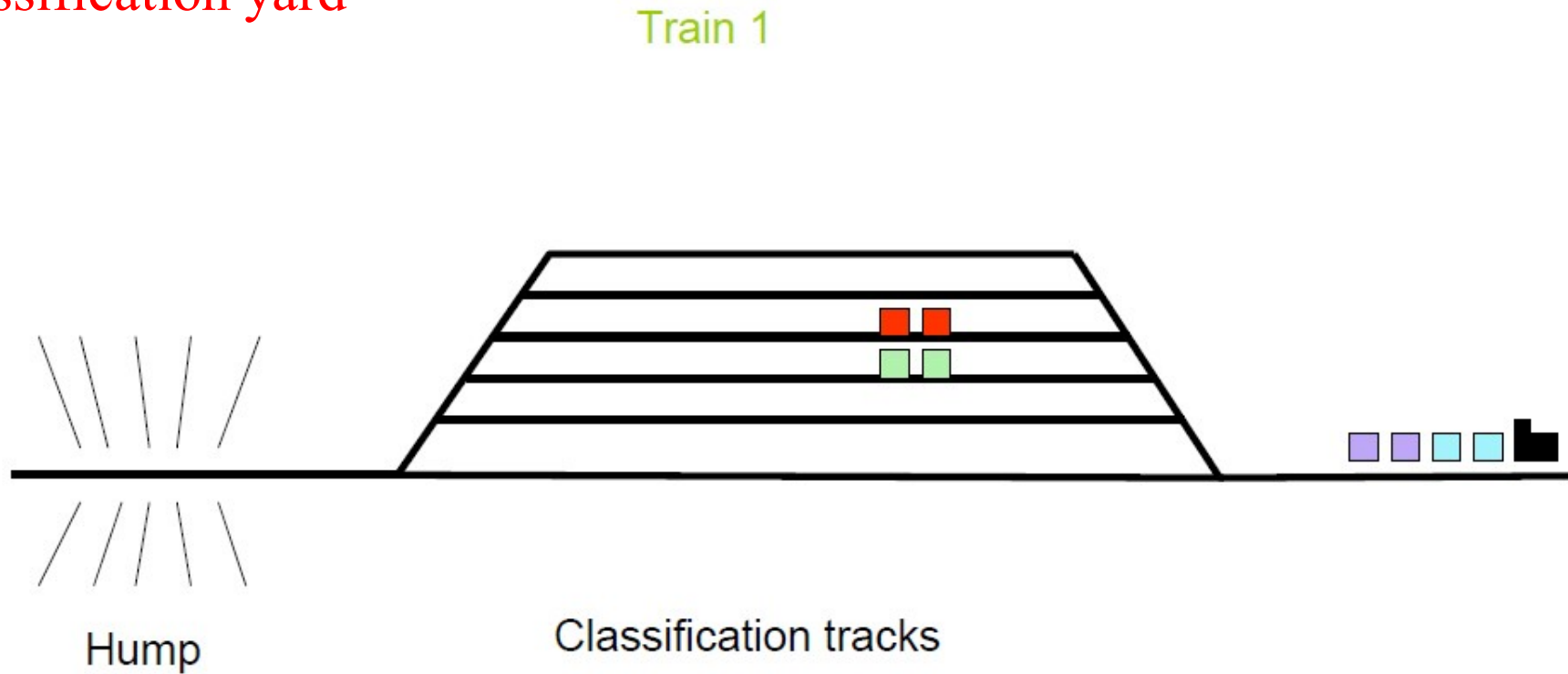
Classification yard



Classification yard



Classification yard

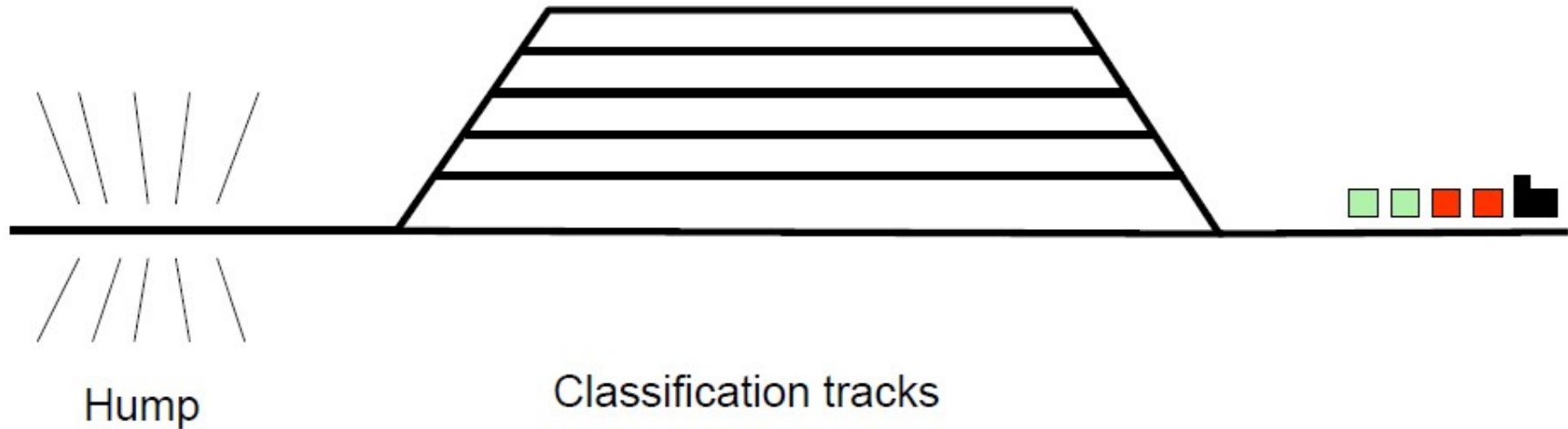


This terminal is as “origin” of the blocks assigned to Train 1 & Train 2



Classification yard

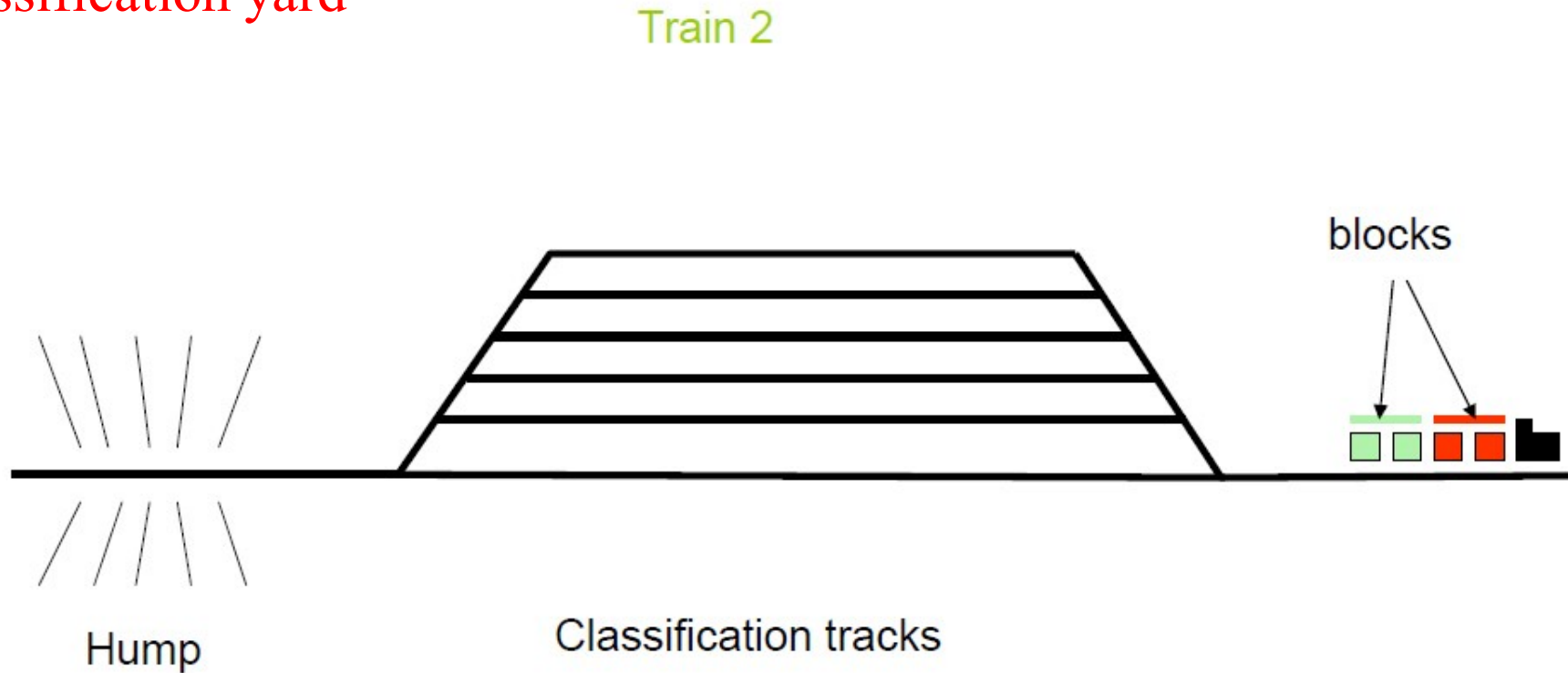
Train 2



This terminal is as “origin” of the blocks assigned to Train 1 & Train 2



Classification yard



This terminal is as “origin” of the blocks assigned to Train 1 & Train 2



Classification yard

- ✓ Number of available classification tracks are not always enough
- ✓ Classification is time-consuming and costly
- ✓ So, we should optimize (Minimize) number of classifications

From Origin to Destination...

- ✓ Once shipments have left their origin and are in their way to their final destination, they can pass through numerous classification yards.
- ✓ In each of these classification yards, shipments may or may not be reclassified again, as befitting the specific circumstances.
- ✓ Every time a shipment is classified, the railway incurs additional costs due to human labor, the use of yard resources, time delays, etc.



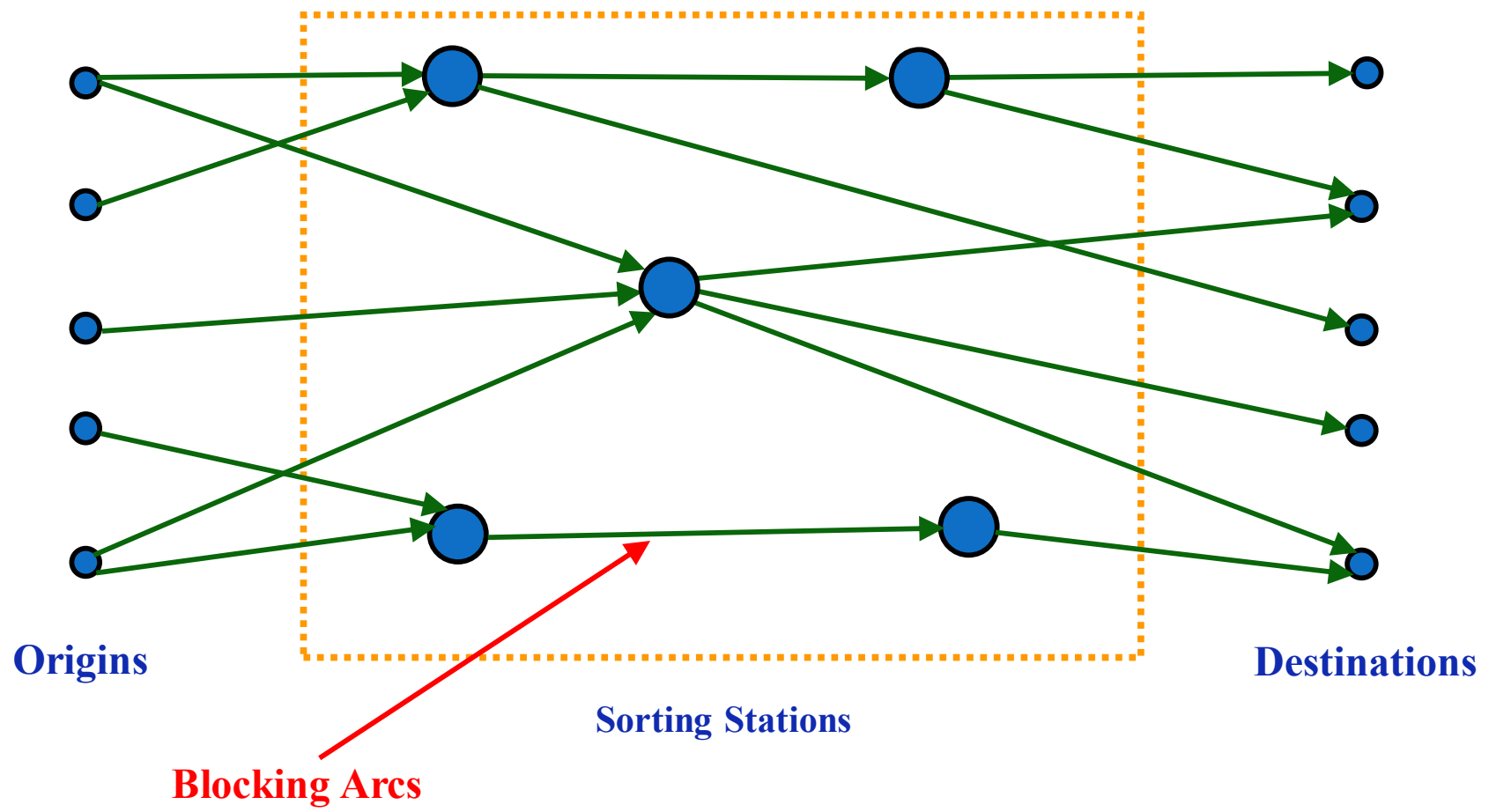
- ✓ In order to **avoid extra costs** (handling costs), commodities that share a final destination (or commodities that have the possibility of traveling together for a portion of their journeys even though they may be destined to different final destinations) are **grouped together** to create a **block**.
- ✓ A block is now paired with a **new origin-destination** set of yards, which may or may not correspond with the origin and destinations of any of the shipments.
- ✓ A train can carry 1 or more blocks.
- ✓ Block A with origin i and destination j . k is a terminal among i and j .



Question: Can the train carrying block A be stopped in Terminal k ?

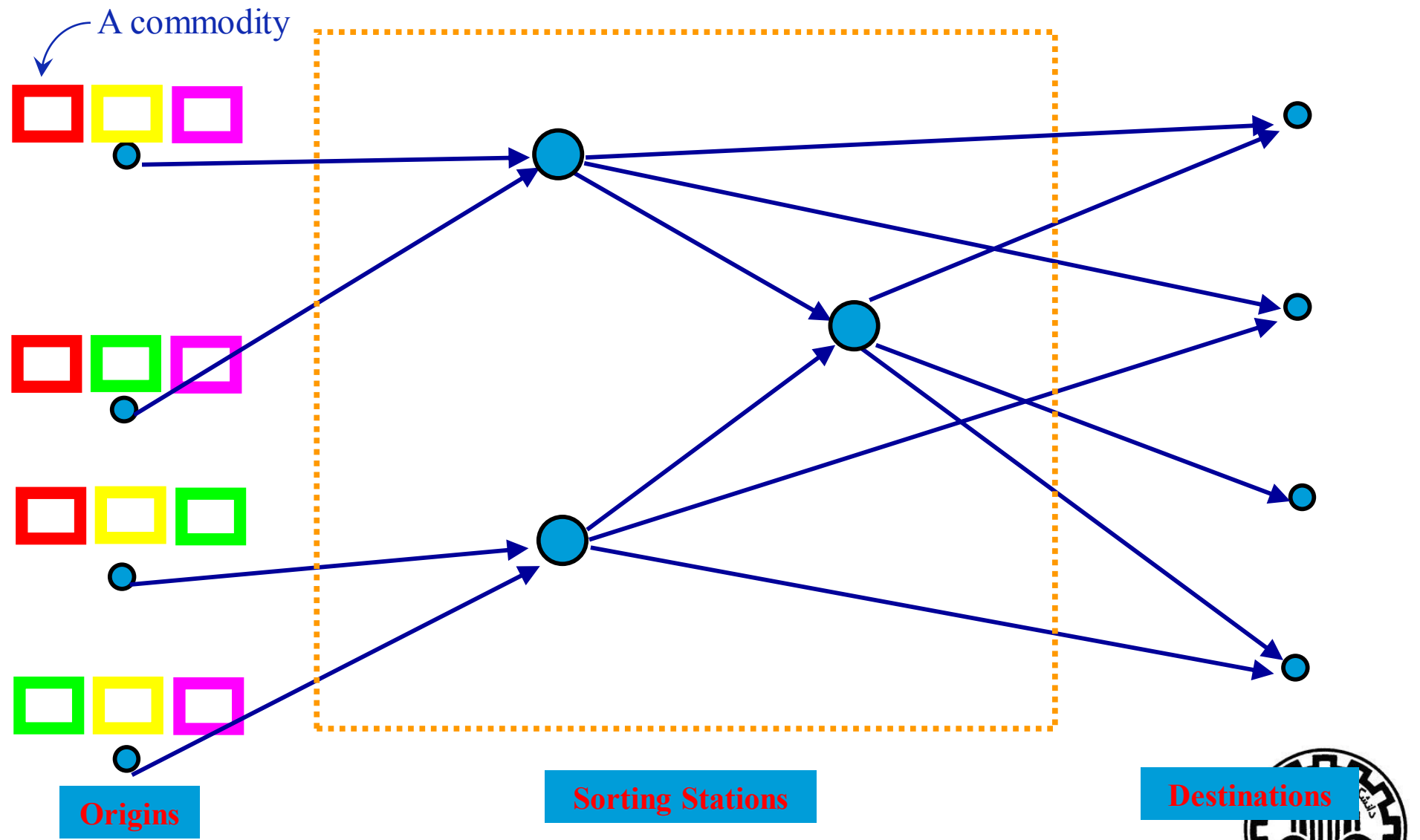
Answer: Yes, But block A must not be reclassified.





Description

Railroad Blocking Problem (RBP)



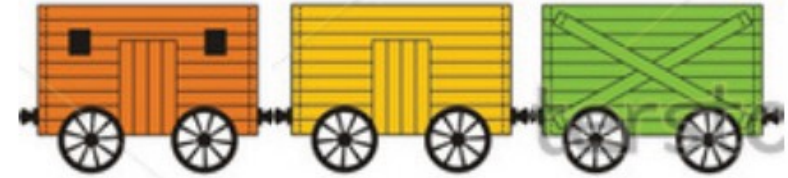
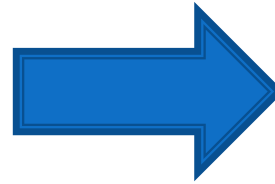
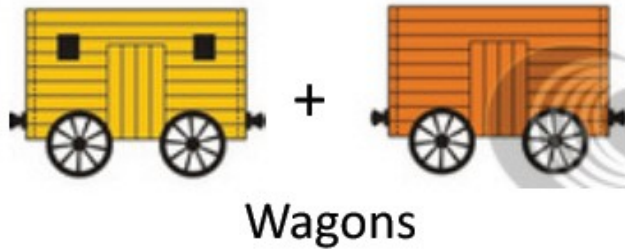
Blocking Problem: Designing the sorting network and route all commodities in it to minimize the weighted sum of travel times and sorting costs

Blocking Plan: determines how to aggregate a large number of commodities into blocks of commodities as they travel from origins to destinations.

- ✓ **Blocking Problem:** developing a plan that describes:
 - which blocks should be assembled at each classification yard?
 - which shipments should be assigned to those blocks?



- ✓ The shipments can pass through a number of classification yards without being reclassified, only getting classified again when they reach destination of block.



Blocks: **same O-D**
Cannot be separated till the block reaches destination

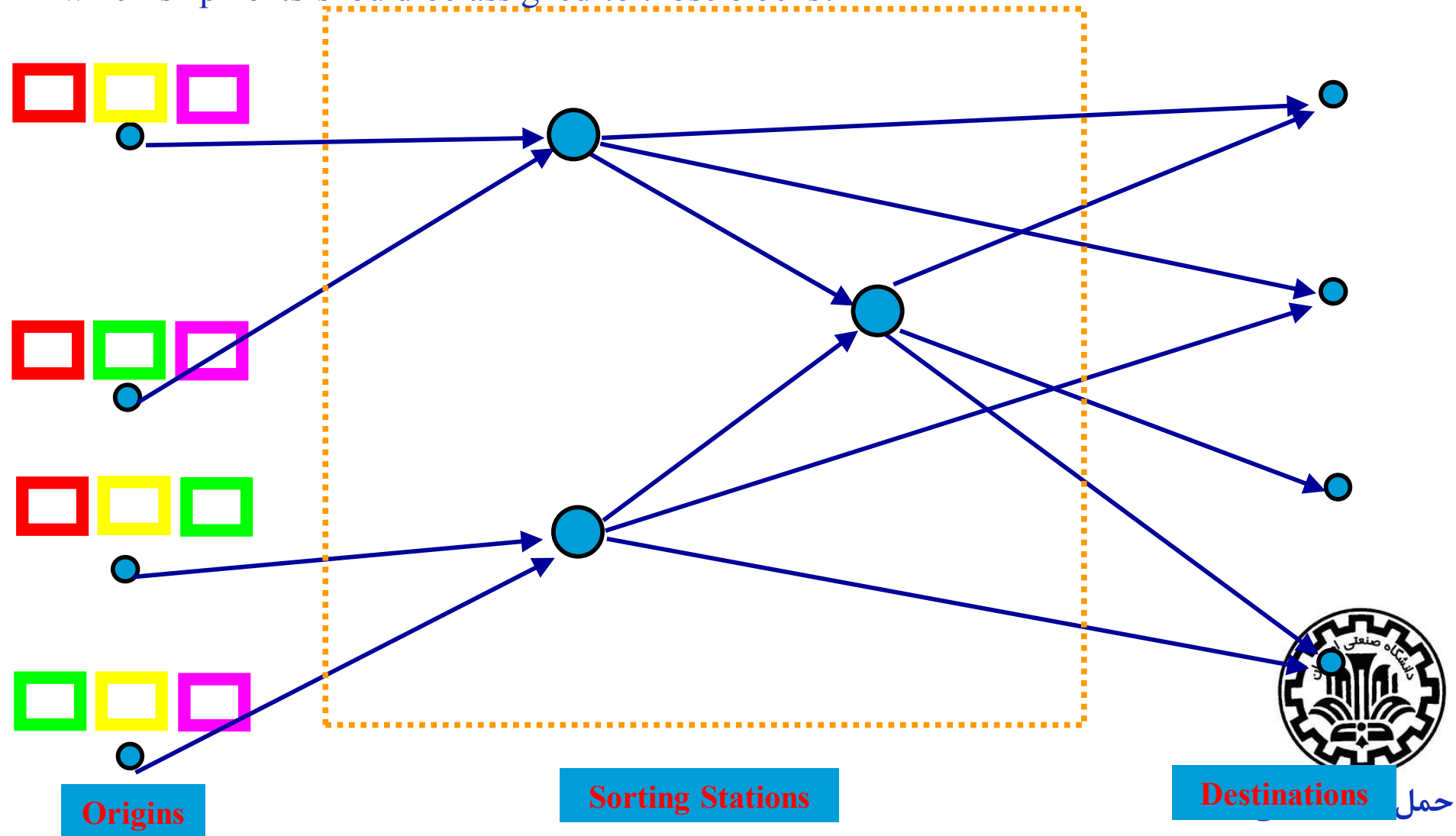
- ✓ Once a block reaches its destination, it is disassembled and the shipments that are not still at their own final destination are assembled into new blocks and continue on their way to their own final destinations
- ✓ Once shipments are grouped together in a block, they get **assigned to trains** consisting of multiple blocks that **share relevant portions of their routes**.

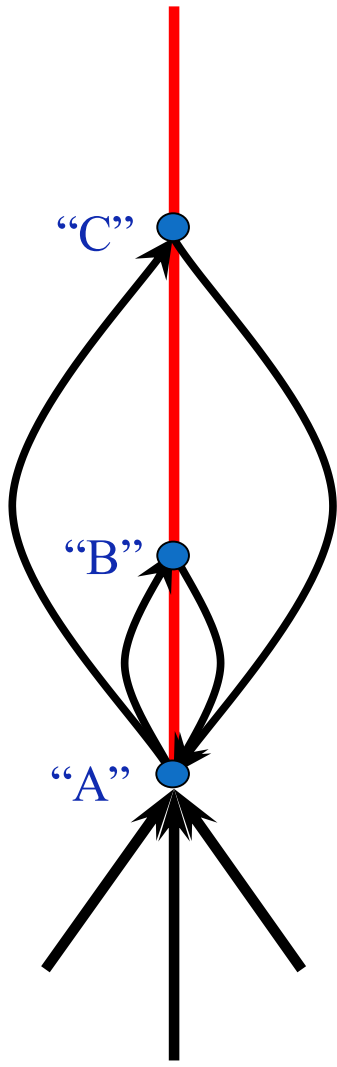


Description

Railroad Blocking Problem (RBP)

- ✓ **Blocking Problem:** developing a plan that describes:
 - which blocks should be assembled at each classification yard?
 - which shipments should be assigned to those blocks?





Block Definition

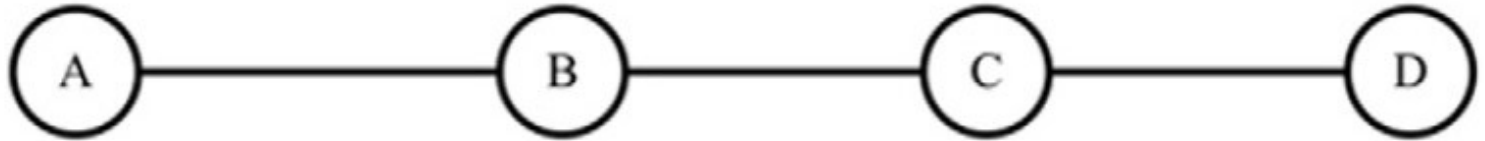
- From "A"
- To "C"

Block Definition

- From "A"
- To "B"

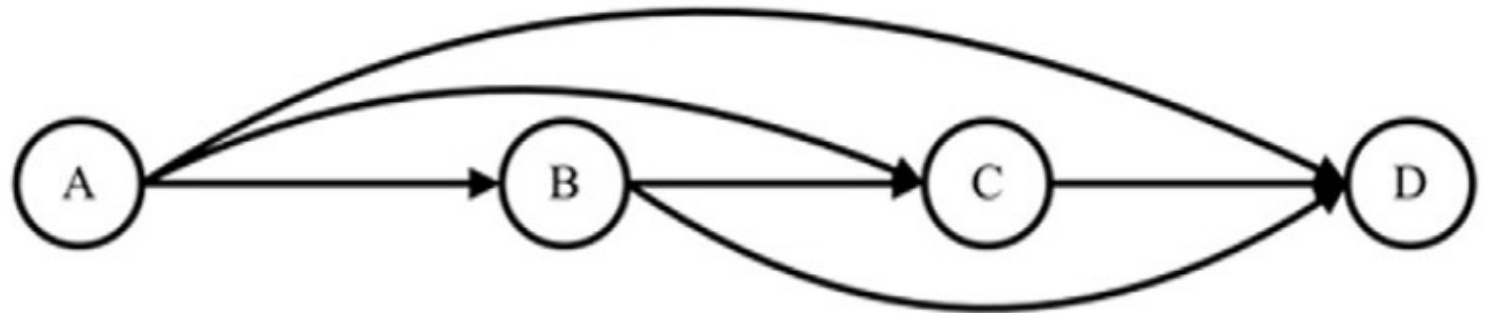


Physical Network:



✓ Physical rail network: the **real** network, consisting of railroad terminals and tracks.

Blocking Network:



✓ The blocking network: a **virtual** network that is overlaid on the physical network.

✓ The blocks are **virtual arcs** which a commodity may take to have uninterrupted service between two terminals that are **not necessarily connected** by a physical link



RBP:

- **Given:**
 - A set of commodities (shipments) with different origins/destinations
- **Determine:**
 - Design the blocking network and route all commodities over blocking network
- **Objective Function:**
 - Minimize the weighted sum of cost of arc travellings and node handlings



RBP:

➤ Constraints:

- Blocking capacities of nodes:

Each block built at any node **needs separate tracks**.

However, because the number of tracks at any node is limited, we can build only a specified number of blocks (blocking capacity of that node).

- Shipment capacities of nodes:

Each node of the network has a limited capacity for the number of cars that can pass through it. Sending more cars than the capacity of the nodes can handle creates **congestion** and may ultimately lead to the complete **breakdown** of operations.

- Flow capacities of blocks:

Each block built in the network may have a specified capacity in terms of the number of cars that can flow on.

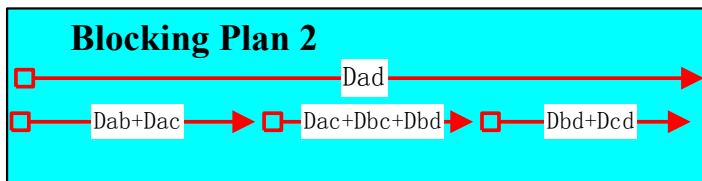
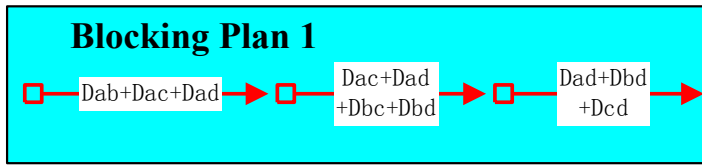
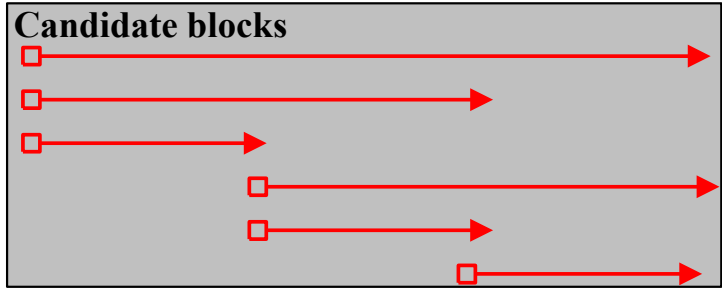
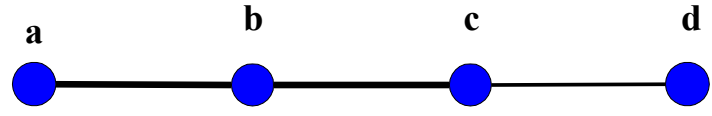
However, each block should carry at least a specified number of cars.



Example 1

Railroad Blocking Problem (RBP)

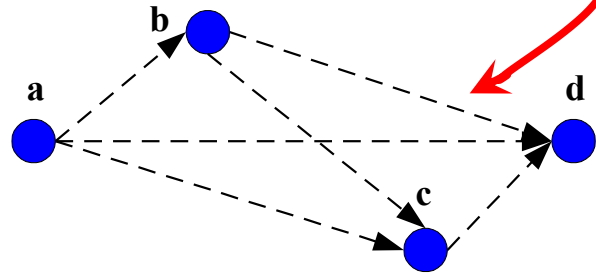
origin \ destination	b	c	d
a	100	100	500
b		150	200
c			50



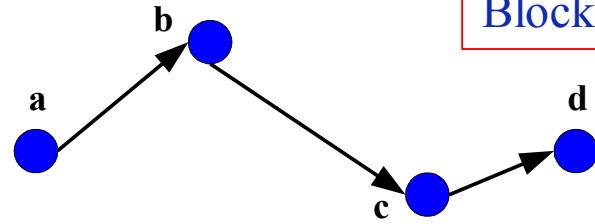
● Station
 □ → Block

Physical network

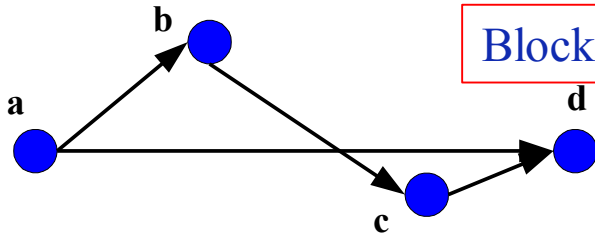
Blocking network



Blocking Plan 1

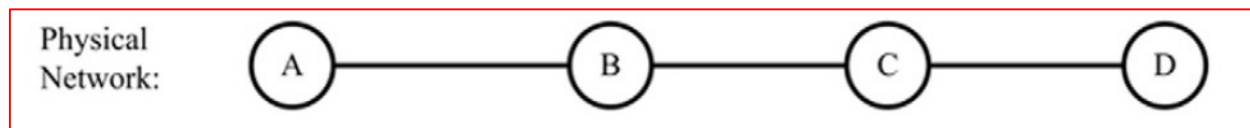


Blocking Plan 2



Example 2

Railroad Blocking Problem (RBP)



✓ commodities (origin-destination pairs of terminals):

A → B with 100 cars

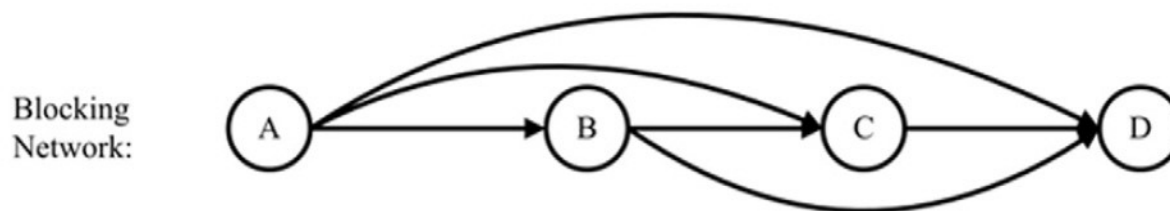
A → C with 80 cars

A → D with 90 cars

Terminals	Shipment capacity	Blocking capacity
Terminal A	270 cars	2 blocks
Terminal B	90 cars	1 block
Terminal C	90 cars	1 block

Shipment capacity constraints: upper bounds on the flow through each node.

Blocking capacity constraints: limits on the out-degree of each node.



Example 2

Railroad Blocking Problem (RBP)

Plan 1



A → B with 100 cars
A → C with 80 cars
A → D with 90 cars

- ✓ Short blocking strategy is used.
- ✓ At terminal A, all three commodities are sorted to move to terminal B.
- ✓ At terminal B, commodity A → B (100 cars) has reached its destination and **leaves** the system by separating them using **classification operation**.
- ✓ Commodities A → C (80 cars) and A → D (90 cars) are blocked by blocking operation to be moved to terminal C.
- ✓ At terminal C, commodity A → C has reached its destination and **leaves** the system by another classification operation
- ✓ commodity A → D is blocked to move to terminal D.
- ✓ Using this blocking plan:
 - 100 cars for commodity A → B use 1 block (A, B)
 - 80 cars for commodity A → C use 2 blocks (A, B) and (B, C)
 - 90 cars for commodity A → D use 3 blocks (A, B), (B, C) and (C, D).
- ✓ Block (A,B) including: A → B, A → C and A → D [270 car classification in terminal A]
- ✓ Block (B,C) including: A → C and A → D [170 car classification in terminal B]
- ✓ Block (C,D) including: A → D [90 car classification in terminal C]
- ✓ Totally, **530** car classification operations.

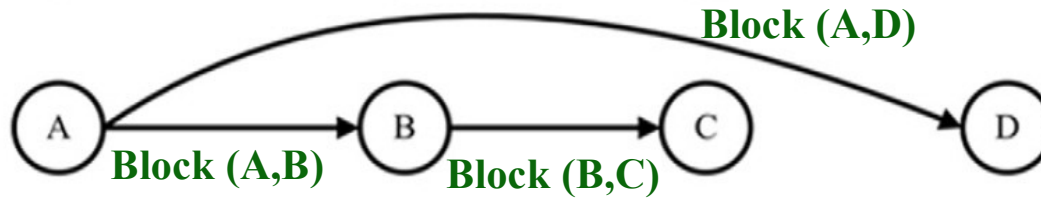
Plan 1 is infeasible (Shipment capacity in Terminal B)



Example 2

Railroad Blocking Problem (RBP)

Plan 2



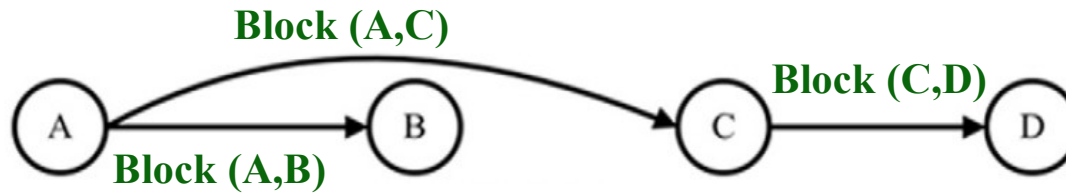
A → B with 100 cars
A → C with 80 cars
A → D with 90 cars

- ✓ Providing a direct block from terminal A to terminal D
- ✓ Cars which travel in block (A,D) still move along physical track through terminals B and C
- ✓ the 90 cars of Commodity A → D now require only one block, (A, D).
- ✓ Totally, **350** car classification operations.
- ✓ Plan 2 is feasible



Example 2

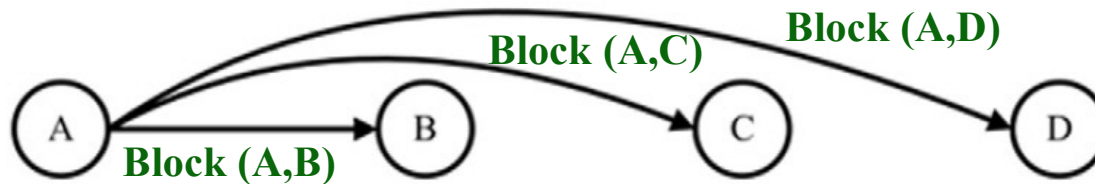
Railroad Blocking Problem (RBP)



Plan 3

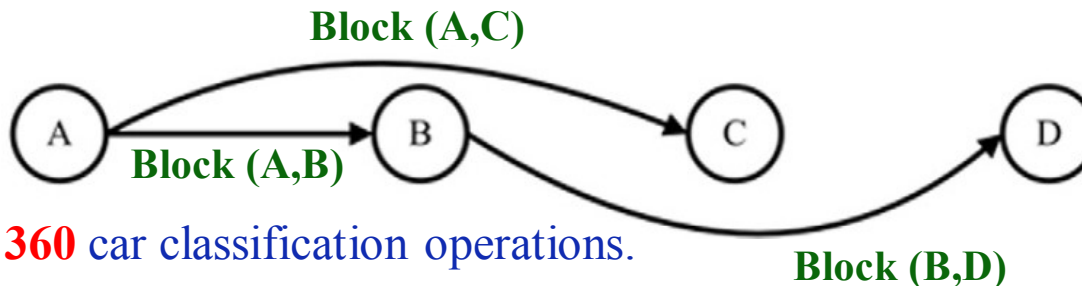
A → B with 100 cars
A → C with 80 cars
A → D with 90 cars

- ✓ Totally, **360** car classification operations.
- ✓ Plan 3 is feasible



Plan 4

- ✓ Totally, **270** car classification operations.
- ✓ Plan 4 is infeasible



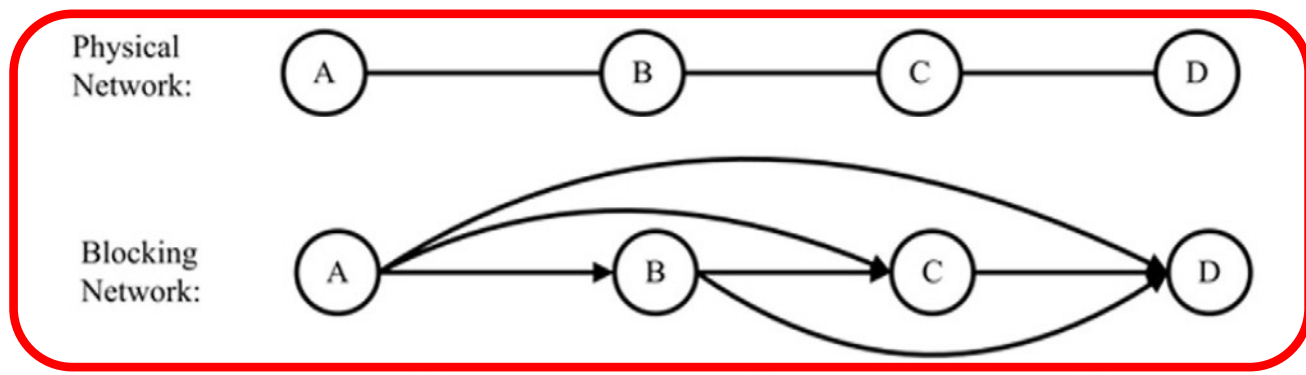
Plan 5

- ✓ Totally, **360** car classification operations.
- ✓ Plan 5 is feasible

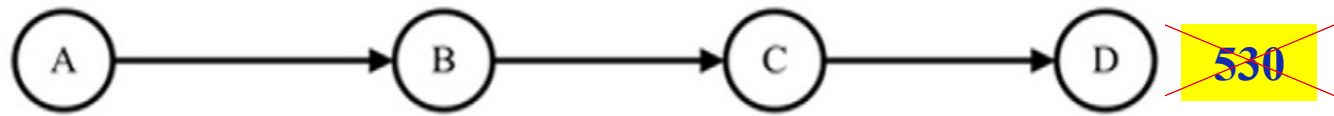


Example 2

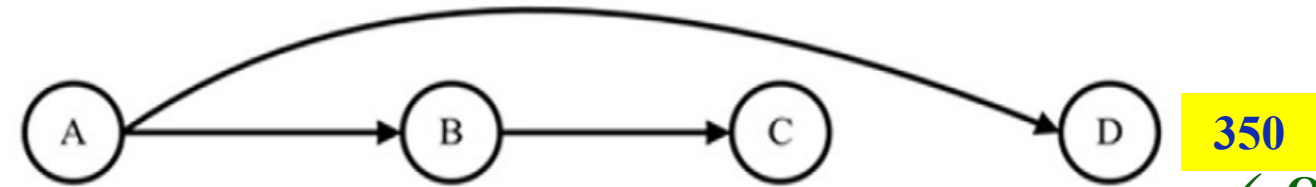
Railroad Blocking Problem (RBP)



Plan 1:

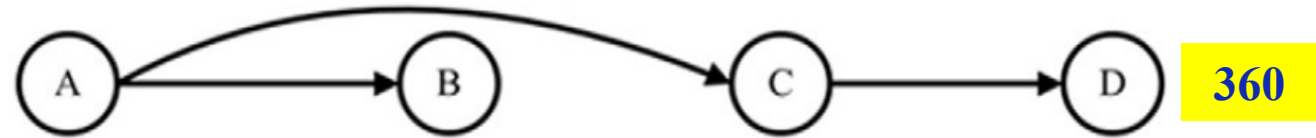


Plan 2:



✓ Optimal Solution

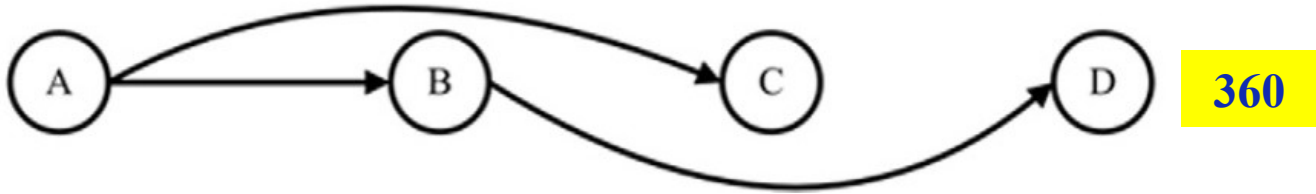
Plan 3:



Plan 4:



Plan 5:



To solve RBP, we must decide which blocks to include in the blocking plan and which blocks to use to deliver each commodity

Decision Variables

- Blocking decisions:

What is the blocking network?

(i.e., how many blocks are made, and what is the origin-destination of each block?)

- Shipment-block sequencing decisions:

How is each origin-destination shipment routed over the blocking network?



“Node-Arc Formulation”

Parameters:

$G = (N, A)$ is the graph with terminal set N and potential blocks set A .

K is the set of all commodities k designated by an origin-destination pair of nodes.

U_k is the volume of commodity k .

$\text{orig}(k)$ is the origin terminal for commodity k .

$\text{dest}(k)$ is the destination terminal for commodity k .

$\text{orig}(a)$ is the origin of potential block a .

$\text{dest}(a)$ is the destination of potential block a .

u_a is the capacity of potential block a .

c_a is the per unit cost of flow on arc a (assumed equal for all commodities).

$B(i)$ is the number of blocks which may be originated at terminal i .

$V(i)$ is the volume which may be classified at terminal i .

Variables:

x_a^k 1 if commodity k is flowing on block a , 0 otherwise.

y_a 1 if block a is included in the blocking network, 0 otherwise.

Ref.: Harry N. Newton, 1996 “network design under budget constraints with application to the railroad blocking problem”, PhD Thesis, Auburn University.

Objective Function:

The objective is to minimize the sum of the costs of delivering each commodity using the blocking network formed by blocks for which $y_a = 1$.

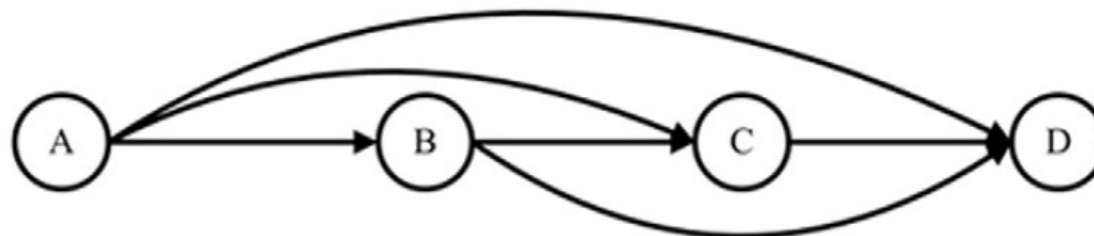
$$\text{Min } Z = \sum_{k \in K} \sum_{a \in A} c_a \cdot v_k \cdot x_a^k$$

the per unit cost of flow on block a

volume of commodity k

1 if commodity k is flowing on block a,
0 otherwise

Blocking Network:



Constraints:

For each terminal, balance equations for the flow of each commodity:

$$\sum_{\substack{a \in A \\ \text{orig}(a)=i}} x_a^k - \sum_{\substack{a \in A \\ \text{dest}(a)=i}} x_a^k = \begin{cases} 1 & \text{orig}(a) = i, \\ -1 & \text{dest}(a) = i, \\ 0 & \text{otherwise,} \end{cases}$$

1 if commodity k
is flowing on block a,
0 otherwise

$$\forall i \in N, k \in K,$$

terminal i

Commodity k



Constraints:

For each potential block:

- 1) flow on blocks which are not built, must be prevented.
- 2) upper bound u_a on flow for blocks -which are built- is enforced.

$$\sum_{k \in K} v^k x_a^k \leq u_a y_a, \quad \forall a \in A,$$

volume of commodity k

capacity of potential block a

Block a

1 if commodity k is flowing on block a,
0 otherwise

1 if block a is included in the blocking network,
0 otherwise



Constraints:

enforcing terminal limit $B(i)$ for sum of the blocks which leave the terminal.

$$\sum_{\substack{a \in A \\ \text{orig}(a)=i}} y_a \leq B(i), \quad \forall i \in N,$$

1 if block a is included
in the blocking network,
0 otherwise

Blocking capacity in terminal i
(Maximum number of blocks
which may be originated at
terminal i).



Constraints:

modelling the volume of cars, which may be classified at each terminal

$$\sum_{k \in K} \sum_{\substack{a \in A \\ \text{orig}(a)=i}} v_k \chi_a^k \leq V(i), \quad \forall i \in N,$$

volume of commodity k

1 if commodity k
is flowing on block a,
0 otherwise

Shipment capacity
(Maximum volume of
commodities which may
be classified at terminal i)

terminal i



Constraints:

$$y_a \in \{0, 1\}, \quad \forall a \in A,$$

$$x_a^k \in \{0, 1\}, \quad \forall a \in A, \quad \forall k \in K.$$



$$\text{Min } Z = \sum_{k \in K} \sum_{a \in A} c_a v_k x_a^k,$$

Subject to:

$$\sum_{\substack{a \in A \\ \text{orig}(a)=i}} x_a^k - \sum_{\substack{a \in A \\ \text{dest}(a)=i}} x_a^k = \begin{cases} 1 & \text{orig}(a) = i, \\ -1 & \text{dest}(a) = i, \\ 0 & \text{otherwise,} \end{cases} \quad \forall i \in N, k \in K,$$

$$\sum_{k \in K} v_k x_a^k \leq u_a y_a, \quad \forall a \in A,$$

$$\sum_{\substack{a \in A \\ \text{orig}(a)=i}} y_a \leq B(i), \quad \forall i \in N,$$

$$\sum_{k \in K} \sum_{\substack{a \in A \\ \text{orig}(a)=i}} v_k x_a^k \leq V(i), \quad \forall i \in N,$$

$$y_a \in \{0, 1\}, \quad \forall a \in A,$$

$$x_a^k \in \{0, 1\}, \quad \forall a \in A, \forall k \in K.$$



A path-based NDP (Network Design Problem) formulation of the railroad blocking problem:

$$(\mathcal{P}) \quad \min \sum_{k \in \mathcal{K}} \sum_{q \in \mathcal{Q}(k)} PC_q^k v^k f_q^k, \quad (1)$$

s.t.

$$\sum_{q \in \mathcal{Q}(k)} \delta_a^q f_q^k \leq y_a, \quad \forall k \in \mathcal{K}, \forall a \in \mathcal{A}, \quad (2)$$

$$\sum_{q \in \mathcal{Q}(k)} f_q^k = 1, \quad \forall k \in \mathcal{K}, \quad (3)$$

$$\sum_{a \in \mathcal{A}} \xi_i^a y_a \leq B(i), \quad \forall i \in \mathcal{N}, \quad (4)$$

$$\sum_{k \in \mathcal{K}} \sum_{q \in \mathcal{Q}(k)} \sum_{a \in \mathcal{A}} v^k \delta_a^q \xi_i^a f_q^k \leq V(i), \quad \forall i \in \mathcal{N}, \quad (5)$$

$$f_q^k \geq 0, \quad \forall q \in \mathcal{Q}(k), \forall k \in \mathcal{K}, \quad (6)$$

$$y_a \in \{0, 1\}, \quad \forall a \in \mathcal{A}, \quad (7)$$

where $PC_q^k = \sum_{a \in \mathcal{A}} C_a \delta_a^q$.

Ref.: C. Barnhart , H. Jin, P.H. Vance, "Railroad Blocking: A Network Design Application", operations research, vol. 48, no. 4, july–august 2000, pp. 603–614

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Railroad Blocking Problem (RBP)

